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The Deductibility of the Aristotelian Modal Syllogism □IAI-4 from the Perspective of Natural Language Information Processing

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Abstract

Aristotelian modal syllogisms characterize the semantic and reasoning properties of Aristotelian quantifiers and modalities. In order to give a consistent explanation for Aristotelian modal syllogisms, this paper reveals the reduction between modal syllogisms on the basis of generalized quantifier theory, set theory, first-order logic, and modern modal logic. To be more specific, this paper firstly proves the validity of the modal syllogism $\Box IAI-4$ based on the truth value definitions of modal categorical propositions, and secondly deduces the other 32 valid modal syllogisms from this syllogism based on related definitions, facts and inference rules. That is to say that there is reducibility among modal syllogisms with different figures and forms. This formal study not only conforms with the needs for formal transformation of all kinds of information in the era of artificial intelligence, but also provides other types of syllogisms with unified mathematical paradigm.

Key words: Aristotelian modal syllogisms; Aristotelian quantifiers; possible worlds; symmetry

1. Introduction

There are various types of syllogisms in natural language, such as Aristotelian syllogisms (Patzig, 1969; Long, 2023; Hui, 2023), Aristotelian modal syllogisms (Łukasiewicz, 1957; Cheng, 2023), and generalized syllogisms (Xiaojun and Baoxiang, 2021), and so on. This paper mainly discusses Aristotelian modal syllogisms.

Aristotelian modal syllogisms have been studied by many scholars, for example, Xiaojun (2020a, 2020b) and Cheng (2023) provide a formal study of Aristotelian modal syllogisms from the perspective of modern logic. Protin (2022) proposes a new deductive system to explain the validity of Aristotelian modal syllogisms. However, many scholars agree that the existing research can't give a consistent explanation of Aristotelian modal syllogisms.

This paper attempts to provide a consistent explanation of Aristotelian modal syllogisms. To this end, on the basis of relevant definitions, facts, and reasoning rules, this paper first proves the validity of the modal syllogism $\Box IAI-4$, and then deduces other 32 valid syllogisms from the modal syllogism $\Box IAI-4$.

2. Relevant Basic Knowledge

Aristotelian syllogisms characterize the semantic and reasoning properties of the following Aristotelian quantifiers: *all, no, some,* and *not all.* An Aristotelian modal syllogism can be obtained by adding at least one possible modality (\diamondsuit) or necessary modality (\Box) to an Aristotelian syllogism. Aristotelian modal syllogisms describe the semantic and reasoning properties of Aristotelian quantifiers and modalities.

In this paper, w, v and z represents the lexical variables of categorical propositions, which are elements in the set W, V and Z, respectively. Let m, n, s and t be propositional variables. The symbol $=_{def}$ means that left can be defined by right.

Aristotelian syllogisms consists of the following: Propositions A, E, I and O (Cheng and Xiaojun, 2023). Proposition A represents 'all ws are vs', which is denoted as all(w, v). Proposition E represents 'no ws are vs', denoted as no(w, v). Proposition I 'some ws are vs', as

some(w, v). The proposition 'not all ws are vs', as not all(w, v).

Definition 1 (truth value of categorical propositions):

(1.1) $all(w, v) =_{def} W \subseteq V;$ (1.2) $some(w, v) =_{def} W \cap V \neq \emptyset;$

(1.3) $no(w, v) =_{def} W \cap V = \emptyset;$ (1.4) $not all(w, v) =_{def} W \not\subseteq V.$

Definition 2 (truth values of modal propositions):

(2.1) $\Box m$ is true, when and only when *m* is true in any possible world ω ;

(2.2) $\diamond m$ is true, when and only when there is at least one possible world ω in which *m* is true.

Definition 3 (inner negation): $Q \neg (w, v) =_{def} Q(w, D-v)$.

Definition 4 (outer negation): $\neg Q(w, v) =_{def} It$ is not that Q(w, v).

According to modal logic (Chagrov and Zakharyaschev, 1997) and generalized quantifier theory (Peters and Westerståhl, 2006), the following facts are provable:

Fact 1 (a necessary proposition implies an assertion one):

 $(1.1) \vdash \Box all(w, v) \rightarrow all(w, v)$, abbreviated as: $\Box A \rightarrow A$;

(1.2) $\vdash \Box no(w, v) \rightarrow no(w, v)$, abbreviated as: $\Box E \rightarrow E$;

(1.3) $\vdash \Box some(w, v) \rightarrow some(w, v)$, abbreviated as: $\Box I \rightarrow I$;

 $(1.4) \vdash \Box not all(w, v) \rightarrow not all(w, v)$, abbreviated as: $\Box O \rightarrow O$.

Fact 2 (an universal proposition implies a particular one):

(2.1) $\vdash all(w, v) \rightarrow some(w, v)$, abbreviated as: $A \rightarrow I$;

(2.2) \vdash *no*(*w*, *v*) \rightarrow *not all*(*w*, *v*), abbreviated as: $E \rightarrow O$;

 $(2.3) \vdash \Box all(w, v) \rightarrow \Box some(w, v)$, abbreviated as: $\Box A \rightarrow \Box I$;

- $(2.4) \vdash \Box no(w, v) \rightarrow \Box not all(w, v)$, abbreviated as: $\Box E \rightarrow \Box O$;
- $(2.5) \vdash \Diamond all(w, v) \rightarrow \Diamond some(w, v)$, abbreviated as: $\Diamond A \rightarrow \Diamond I$;
- $(2.6) \vdash \Diamond no(w, v) \rightarrow \Diamond not \ all(w, v)$, abbreviated as: $\Diamond E \rightarrow \Diamond O$.

Fact 3 (symmetry):

 $(3.1) \vdash some(w, v) \leftrightarrow some(v, w); \qquad (3.2) \vdash \Box some(w, v) \leftrightarrow \Box some(v, w);$

 $(3.3) \vdash \diamond some(w, v) \leftrightarrow \diamond some(v, w);$ $(3.4) \vdash no(w, v) \leftrightarrow no(v, w);$ $(3.6) \vdash \Diamond no(w, v) \leftrightarrow \Diamond no(v, w).$ $(3.5) \vdash \Box no(w, v) \leftrightarrow \Box no(v, w);$ Fact 4 (inner negation): $(4.1) \vdash all(w, v) \leftrightarrow no \neg (w, v);$ $(4.2) \vdash no(w, v) \leftrightarrow all \neg (w, v);$ $(4.3) \vdash some(w, v) \leftrightarrow not all \neg (w, v);$ $(4.4) \vdash not all(w, v) \leftrightarrow some_{\neg}(w, v).$ Fact 5 (outer negation): $(5.1) \vdash \neg not all(w, v) \leftrightarrow all(w, v);$ $(5.2) \vdash \neg all(w, v) \leftrightarrow not all(w, v);$ $(5.3) \vdash \neg no(w, v) \leftrightarrow some(w, v);$ $(5.4) \vdash \neg some(w, v) \leftrightarrow no(w, v).$ **Fact 6** (dual): (6.1) $\vdash \neg \Box O(w, v) \leftrightarrow \Diamond \neg O(w, v)$; $(6.2) \vdash \neg \Diamond O(w, v) \leftrightarrow \Box \neg O(w, v)$ *v*).

Rule 1 (subsequent weakening): If $\vdash (m \land n \rightarrow s)$ and $\vdash (s \rightarrow t)$, then $\vdash (m \land n \rightarrow t)$.

Rule 2 (anti-syllogism): If $\vdash (m \land n \rightarrow s)$, then $\vdash (\neg s \land m \rightarrow \neg n)$ or $\vdash (\neg s \land n \rightarrow \neg m)$.

3. The Validity of the Syllogism $\Box IAI-4$

In order to discuss the reducibility of modal syllogisms based on the syllogism $\Box IAI-4$, it is important to prove the validity of the syllogism $\Box IAI-4$.

Theorem 1 (\Box *IAI-4*): \Box *some*(v, z) \land *all*(z, w) \rightarrow *some*(w, v) is valid.

Proof: The modal syllogism \Box some $(v, z) \land all(z, w) \rightarrow some(w, v)$ can be abbreviated as \Box *IAI-4*. Suppose that \Box some(v, z) and all(z, w) are true, then some(v, z) is true in any possible world according to the Definition (2.1). Due to that any real world is a possible world, some(v, z) is true in any real world, thus $V \cap Z \neq \emptyset$ is true in line with Definition (1.2). And all(z, w) is true in real world, then $Z \subseteq W$ is true according to the Definition (1.1). Now it follows $W \cap V \neq \emptyset$ is true in any real world. Hence some(w, v) is true according to the Definition (1.2). The above proves that the syllogism $\Box some(v, z) \land all(z, w) \rightarrow some(w, v)$ is valid.

4. The Other 32 Modal Syllogisms Derived from DIAI-4

According to Theorem 1, syllogism $\Box IAI-4$ is valid. '(1) $\Box IAI-4 \rightarrow \Box IAI-3$ ' in Theorem 2 means that the syllogism $\Box IAI-3$ can be derived from the syllogism $\Box IAI-4$. That is to say that there is reducibility between these two syllogisms. The others are similar.

Theorem 2: The following 32 valid modal syllogisms can be inferred from $\Box IAI-4$:

- (1) $\Box IAI-4 \rightarrow \Box IAI-3$
- (2) $\Box IAI-4 \rightarrow A \Box II-1$
- $(3) \Box IAI-4 \rightarrow A \Box II-3$
- (4) $\Box IAI-4 \rightarrow E \Box IO-4$
- (5) $\Box IAI-4 \rightarrow AE \diamondsuit E-4$
- (6) $\Box IAI-4 \rightarrow \Box IAI-3 \rightarrow E \Box IO-2$
- (7) $\Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamondsuit E-1$
- $(8) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow \Box OAO-3$
- $(9) \Box IAI-4 \rightarrow A \Box II-1 \rightarrow E \Box IO-3$
- (10) $\Box IAI-4 \rightarrow A \Box II-1 \rightarrow AE \diamondsuit E-2$
- $(11) \Box IAI-4 \rightarrow A \Box II-1 \rightarrow E \Box IO-1$
- (12) $\Box IAI-4 \rightarrow A \Box II-3 \rightarrow EA \diamondsuit E-2$
- (13) $\Box IAI-4 \rightarrow AE \diamond E-4 \rightarrow AE \diamond O-4$
- (14) $\Box IAI-4 \rightarrow \Box IAI-3 \rightarrow E \Box IO-2 \rightarrow A \Box OO-2$
- $(15) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamondsuit E-1 \rightarrow EA \diamondsuit O-1$
- (16) $\Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow AA \diamond A-1$
- (17) $\Box IAI-4 \rightarrow A \Box II-1 \rightarrow AE \diamondsuit E-2 \rightarrow AE \diamondsuit O-2$
- (18) $\Box IAI-4 \rightarrow A \Box II-3 \rightarrow EA \diamond E-2 \rightarrow EA \diamond O-2$
- (19) $\Box IAI-4 \rightarrow AE \diamond E-4 \rightarrow AE \diamond O-4 \rightarrow E \Box AO-4$
- $(20) \Box IAI-4 \rightarrow AE \diamondsuit E-4 \rightarrow AE \diamondsuit O-4 \rightarrow \Box AAI-4$
- $(21) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow E \Box AO-2$

- $(22) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow \Box AAI-3$
- $(23) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow AA \diamond I-1$
- $(24) \Box IAI-4 \rightarrow A \Box II-1 \rightarrow AE \diamondsuit E-2 \rightarrow AE \diamondsuit O-2 \rightarrow E \Box AO-3$
- $(25) \Box IAI-4 \rightarrow A \Box II-1 \rightarrow AE \diamondsuit E-2 \rightarrow AE \diamondsuit O-2 \rightarrow A \Box AI-1$
- $(26) \Box IAI-4 \rightarrow A \Box II-3 \rightarrow EA \diamond E-2 \rightarrow EA \diamond O-2 \rightarrow E \Box AO-1$
- $(27) \Box IAI-4 \rightarrow A \Box II-3 \rightarrow EA \diamond E-2 \rightarrow EA \diamond O-2 \rightarrow A \Box AI-3$
- $(28) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow E \Box AO-2 \rightarrow A \Box EO-2$
- $(29) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow \Box AAI-3 \rightarrow \Box EAO-3$
- $(30) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamondsuit E-1 \rightarrow EA \diamondsuit O-1 \rightarrow AA \diamondsuit I-1 \rightarrow AA \diamondsuit I-4$
- $(31) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow E \Box AO-2 \rightarrow A \Box EO-2 \rightarrow A \Box EO-4$
- $(32) \Box IAI-4 \rightarrow \Box IAI-3 \rightarrow EA \diamond E-1 \rightarrow EA \diamond O-1 \rightarrow \Box AAI-3 \rightarrow \Box EAO-3 \rightarrow \Box EAO-4$

Proof:

(i.e. $\Box IAI-4$, Theorem 1) [1] $\vdash \Box some(v, z) \land all(z, w) \rightarrow some(w, v)$ (i.e. $\Box IAI-3$, by [1] and Fact (3.2)) $[2] \vdash \Box some(z, v) \land all(z, w) \rightarrow some(w, v)$ $[3] \vdash \Box some(v, z) \land all(z, w) \rightarrow some(v, w)$ (i.e. $A \Box II-I$, by [1] and Fact (3.1)) $[4] \vdash \Box some(z, v) \land all(z, w) \rightarrow some(v, w)$ (i.e. $A \Box II$ -3, by [1], Fact (3.1) and Fact (3.2)) $[5] \vdash \neg some(w, v) \land \Box some(v, z) \rightarrow \neg all(z, w)$ (by [1] and Rule 2)[6] $\vdash no(w, v) \land \Box some(v, z) \rightarrow not all(z, w)$ (i.e. $E \Box IO-4$, by [5], Fact (5.2) and Fact (5.4)) $[7] \vdash \neg some(w, v) \land all(z, w) \rightarrow \neg \Box some(v, z)$ (by [1] and Rule 2) [8] $\vdash no(w, v) \land all(z, w) \rightarrow \Diamond \neg some(v, z)$ (by [7], Fact (5.4) and Fact (6.1)) $[9] \vdash no(w, v) \land all(z, w) \rightarrow \Diamond no(v, z)$ (i.e. $AE \diamond E-4$, by [8] and Fact (5.4)) (by [2] and Rule 2) $[10] \vdash \neg some(w, v) \land \Box some(z, v) \rightarrow \neg all(z, w)$ $[11] \vdash no(w, v) \land \Box some(z, v) \rightarrow not \ all(z, w)$ (i.e. $E \Box IO-2$, by [10], Fact (5.2) and Fact (5.4)) $[12] \vdash \neg some(w, v) \land all(z, w) \rightarrow \neg \Box some(z, v)$ (by [2] and Rule 2)(by [12], Fact (5.4) and Fact (6.1)) $[13] \vdash no(w, v) \land all(z, w) \rightarrow \Diamond \neg some(z, v)$

- $[14] \vdash no(w, v) \land all(z, w) \rightarrow \Diamond no(z, v)$
- $[15] \vdash \Box not all \neg (z, v) \land all(z, w) \rightarrow not all \neg (w, v)$
- $[16] \vdash \Box not all(z, D-v) \land all(z, w) \rightarrow not all(w, D-v)$
- $[17] \vdash \neg some(v, w) \land \Box some(v, z) \rightarrow \neg all(z, w)$
- $[18] \vdash no(v, w) \land \Box some(v, z) \rightarrow not \ all(z, w)$
- $[19] \vdash \neg some(v, w) \land all(z, w) \rightarrow \neg \Box some(v, z)$
- $[20] \vdash \neg some(v, w) \land all(z, w) \rightarrow \Diamond \neg some(v, z)$
- $[21] \vdash no(v, w) \land all(z, w) \rightarrow \Diamond no(v, z)$
- $[22] \vdash \Box some(v, z) \land no \neg(z, w) \rightarrow not \ all \neg(v, w)$
- $[23] \vdash \Box some(v, z) \land no(z, D-w) \rightarrow not \ all(v, D-w)$
- $[24] \vdash \neg some(v, w) \land all(z, w) \rightarrow \neg \Box some(z, v)$
- $[25] \vdash \neg some(v, w) \land all(z, w) \rightarrow \Diamond \neg some(z, v)$
- $[26] \vdash no(v, w) \land all(z, w) \rightarrow \Diamond no(z, v)$
- $[27] \vdash no(w, v) \land all(z, w) \rightarrow \Diamond not \ all(v, z)$
- $[28] \vdash all \neg (w, v) \land \Box not \ all \neg (z, v) \rightarrow not \ all(z, w)$
- $[29] \vdash all(w, D-v) \land \Box not \ all(z, D-v) \rightarrow not \ all(z, w)$
- $[30] \vdash no(w, v) \land all(z, w) \rightarrow \diamondsuit not \ all(z, v)$
- $[31] \vdash all \neg (w, v) \land all(z, w) \rightarrow \Diamond all \neg (z, v)$
- $[32] \vdash all(w, D-v) \land all(z, w) \rightarrow \diamondsuit all(z, D-v)$
- $[33] \vdash no(v, w) \land all(z, w) \rightarrow \Diamond not \ all(v, z)$
- $[34] \vdash no(v, w) \land all(z, w) \rightarrow \Diamond not \ all(z, v)$
- $[35] \vdash \neg \diamondsuit not all(v, z) \land no(w, v) \rightarrow \neg all(z, w)$
- $[36] \vdash \Box \neg not all(v, z) \land no(w, v) \rightarrow \neg all(z, w)$
- $[37] \vdash \Box all(v, z) \land no(w, v) \rightarrow not \ all(z, w)$
- $[38] \vdash \neg \diamondsuit not all(v, z) \land all(z, w) \rightarrow \neg no(w, v)$

- (i.e. $EA \diamond E-I$, by [13] and Fact (5.4))
 - (by [2] and Fact (4.3))
- (i.e. $\Box OAO-3$, by [15] and Definition 3)
 - (by [3] and Rule 2)
- (i.e. *E* \Box *IO-3*, by [17], Fact (5.2) and Fact (5.4))
 - (by [3] and Rule 2)
 - (by [19] and Fact (6.1))
 - (i.e. $AE \diamond E$ -2, by [20] and Fact (5.4))
 - (by [3], Fact (4.1) and Fact (4.3))
 - (i.e. $E \Box IO-1$, by [22] and Definition 3)
 - (by [4] and Rule 2)
 - (by [24] and Fact (6.1))
 - (i.e. $EA \diamond E-2$, by [25] and Fact (5.4))
 - (i.e. $AE \diamondsuit O-4$, by [9], Fact (2.6) and Rule 1)
 - (by [11], Fact (4.2) and Fact (4.3))
 - (i.e. $A \square OO-2$, by [28] and Definition 3)
 - (i.e. $EA \diamondsuit O-1$, by [14], Fact (2.6) and Rule1)
 - (by [14] and Fact (4.2))
 - (i.e. $AA \diamond A$ -1, by [31] and Definition 3)
 - (i.e. *AE*◊*O*-2, by [21], Fact (2.6) and Rule 1)
 - (i.e. *EA*�*O*-2, by [26], Fact (2.6) and Rule 1)
 - (by [27] and Rule 2)
 - (by [35] and Fact (6.2))
- (i.e. *E* \Box *AO*-4, by [36], Fact (5.1) and Fact (5.2))
 - (by [27] and Rule 2)

- $[39] \vdash \Box \neg not \ all(v, z) \land all(z, w) \rightarrow \neg no(w, v)$
- $[40] \vdash \Box all(v, z) \land all(z, w) \rightarrow some(w, v)$
- $[41] \vdash \neg \diamondsuit not all(z, v) \land no(w, v) \rightarrow \neg all(z, w)$
- $[42] \vdash \Box \neg not \ all(z, v) \land no(w, v) \rightarrow \neg all(z, w)$
- $[43] \vdash \Box all(z, v) \land no(w, v) \rightarrow not \ all(z, w)$
- $[44] \vdash \neg \diamondsuit not all(z, v) \land all(z, w) \rightarrow \neg no(w, v)$
- $[45] \vdash \Box \neg not \ all(z, v) \land all(z, w) \rightarrow \neg no(w, v)$
- $[46] \vdash \Box all(z, v) \land all(z, w) \rightarrow some(w, v)$
- $[47] \vdash all \neg (w, v) \land all(z, w) \rightarrow \diamondsuit some \neg (z, v)$
- $[48] \vdash all(w, D-v) \land all(z, w) \rightarrow \Diamond some(z, D-v)$
- $[49] \vdash \neg \diamondsuit not all(v, z) \land no(v, w) \rightarrow \neg all(z, w)$
- $[50] \vdash \Box \neg not \ all(v, z) \land no(v, w) \rightarrow \neg all(z, w)$
- $[51] \vdash \Box all(v, z) \land no(v, w) \rightarrow not \ all(z, w)$
- $[52] \vdash \neg \diamondsuit not all(v, z) \land all(z, w) \rightarrow \neg no(v, w)$
- $[53] \vdash \Box \neg not all(v, z) \land all(z, w) \rightarrow \neg no(v, w)$
- $[54] \vdash \Box all(v, z) \land all(z, w) \rightarrow some(v, w)$
- $[55] \vdash \neg \diamondsuit not all(z, v) \land no(v, w) \rightarrow \neg all(z, w)$
- $[56] \vdash \Box \neg not \ all(z, v) \land no(v, w) \rightarrow \neg all(z, w)$
- $[57] \vdash \Box all(z, v) \land no(v, w) \rightarrow not \ all(z, w)$
- $[58] \vdash \neg \diamondsuit not all(z, v) \land all(z, w) \rightarrow \neg no(v, w)$
- $[59] \vdash \Box \neg not \ all(z, v) \land all(z, w) \rightarrow \neg no(v, w)$
- $[60] \vdash \Box all(z, v) \land all(z, w) \rightarrow some(v, w)$
- $[61] \vdash \Box no \neg (z, v) \land all \neg (w, v) \rightarrow not \ all(z, w)$
- $[62] \vdash \Box no(z, D-v) \land all(w, D-v) \rightarrow not \ all(z, w)$
- $[63] \vdash \Box no \neg (z, v) \land all(z, w) \rightarrow not \ all \neg (w, v)$

- (by [38] and Fact (6.2))
- (i.e. □*AAI-4*, by [39], Fact (5.1) and Fact (5.3))
 - (by [30] and Rule 2)
 - (by [41] and Fact (6.2))
- (i.e. $E \Box AO-2$, by [42], Fact (5.1) and Fact (5.2))
 - (by [30] and Rule 2)
 - (by [44] and Fact (6.2))
 - (i.e. □*AAI-3*, by [45], Fact (5.1) and Fact (5.3))
 - (by [30], Fact (4.2) and Fact (4.4))
 - (i.e. $AA \diamond I$ -1, by [47] and Definition 3)
 - (by [33] and Rule 2)
 - (by [49] and Fact (6.2))
- (i.e. $E \Box AO-3$, by [50], Fact (5.1) and Fact (5.2))
 - (by [33] and Rule 2)
 - (by [52] and Fact (6.2))
- (i.e. $A \Box AI I$, by [53], Fact (5.1) and Fact (5.3))
 - (by [34] and Rule 2)
 - (by [55] and Fact (6.2))
- (i.e. $E \Box AO-1$, by [56], Fact (5.1) and Fact (5.2))
 - (by [34] and Rule 2)
 - (by [58] and Fact (6.2))
 - (i.e. *A* \Box *AI-3*, by [59], Fact (5.1) and Fact (5.3))
 - (by [43], Fact (4.1) and Fact (4.2))
 - (i.e. $A \square EO-2$, by [61] and Definition 3)
 - (by [46], Fact (4.1) and Fact (4.3))



At this point, the other 32 modal syllogisms have been deduced from the syllogism $\Box IAI-4$.

5. Conclusion

To sum up, in order to give a consistent explanation for Aristotelian modal syllogisms, this paper shows the reduction between modal syllogisms on the basis of generalized quantifier theory, set theory, first-order logic, and modern modal logic. To be more specific, this paper firstly proves the validity of the modal syllogism $\Box IAI-4$ based on the truth value definitions of modal categorical propositions, and secondly deduces the other 32 valid modal syllogisms from this syllogism based on related definitions, facts and inference rules. In fact, one can derive more valid syllogisms when he continues to derive by means of similar reasoning steps. In other words, there is reducibility among modal syllogisms with different figures and forms.

For future study, can we consider using this research method to discuss the validity and reducibility of other types of syllogisms, such as generalized syllogisms and generalized modal syllogisms?

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